

Memorial University of Newfoundland
Engineering 4862 MICROPROCESSORS
Assignment 2

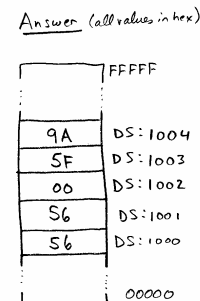
Unless otherwise noted, please show all relevant calculations, and explain your answers where appropriate.

0. List all 8086/8088 registers that can be accessed as both words and bytes.

Accumulator AX	AH	AL
Base BX	BH	BL
Counter CX	CH	CL
Data DX	DH	DL

1. Use a memory map to show the contents of memory locations DS: 1000H to DS: 1004H after all of the following instructions have executed:

	Memory Modification
<i>MOV AX, 56H</i>	None
<i>MOV [1001H], AX</i>	[DS:1001] ← 56H [DS:1002] ← 00H
<i>MOV [1003H], 9A5FH</i>	[DS:1003] ← 5FH [DS:1004] ← 9AH
<i>MOV [1000H], AL</i>	[DS:1000] ← 56H



2. Assume that (CS) = B795H; (DS) = 2000H; (SS) = 0AD4H; (ES) = 30FFH; (SP) = 00FFH; (BP) = 1DF7H; (AX) = 0B24H; (CX) = 1EE4H; (SI) = 3C00H; (DX) = 329FH.

a. Calculate the beginning and ending addresses for all of the segments.

Segment	Start (Base)	End
Code, CS=B795H	B7950H	C794FH
Data, DS=2000H	20000H	2FFFFH
Extra, ES=30FFH	30FF0H	40FEFH
Stack, SS=0AD4H	0AD40H	1AD3FH

b. Suppose that the offset address for the next instruction to be fetched (that is, the contents of IP) is 902DH. Calculate the physical address from where the next instruction will be fetched.

Segment: Code, CS=B795H
PA = CS:IP = B795:9020 = C097DH

c. What will be the contents of BL and AX after the following instruction is executed? Give your results in decimal and hexadecimal.

MOV BL, AL
BL ← AL, AX= 0B24H
So, BL=24H = 36D
AL=24H (unchanged), AX= 0B24H=2852D

d. Calculate the physical addresses of the memory locations referred to in the following instructions and the contents of all the location(s):

AND CX, [1200H]

[1200] is the memory reference, the PA = DS:1200=21200H. Used to access one word.

OR ES:[0E8C9H], SI

ES for Segment override. [0E8C9H] for memory reference. PA = ES:E8C9H = 3F8B9H
Both 3F8B9H and 3F8BAH are referred.

PUSH DX

Memory reference is implicit, as the stack is used.

Stack address = SS:SP = 0AD4H:00FFH = 0AE3FH

The memory accessed is not this location, but at 2 less than this. A word is accessed, the location 0AE3DH and 0AE3EH are used.

3. **What is wrong with, or missing from, each of the following instructions:**

a. **MOV ES, 249EH**

ES is a segment register. It is invalid to assign a segment register directly via an immediate operand.

b. **MOV [BX+3EH], 2** (Hint: the use of PTR directive)

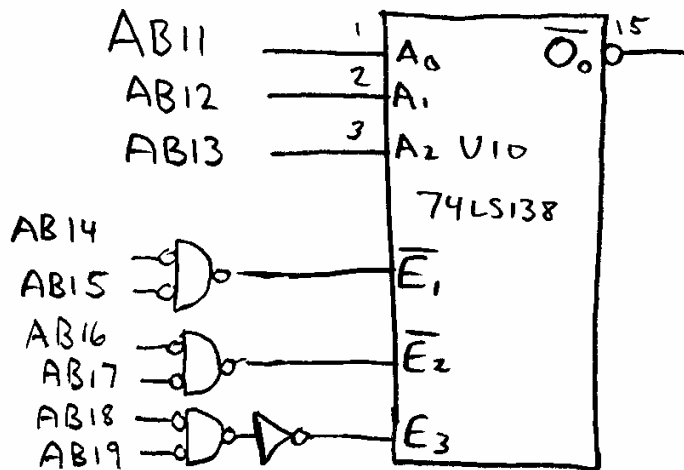
2 can be either a byte or a word operand. The assembler will complain because it does not know how many bytes the value '2' requires. Use BYTE PTR 2 or WORD PTR 2.

c. **MOV [78H], [79H]**

Attempting to have a memory to memory data transfer. Invalid addressing mode.

4. **In lab 1, we found that the MUN-88 single board computer has a number of mirror images. Without replacing the 3x8 decoder or the current 2KB SRAM chip, propose a simple scheme to eliminate these mirror images for the RAM and draw a sketch to show that. (Hint: The mirror images are caused by the don't care lines of the address bus. You can use those lines connect to chip select, enable pins of the decoder or SRAM chips).**

The mirror images are actually caused by the three most significant address lines, which are set to don't cares. One we can give a fixed pattern for those address lines, we will eliminate the confusion. There are many solutions for this question. The diagram below show one possible solution. This requires the addition of one 74LS32 chip (2 input OR gate). The important aspect of this scheme is that AB14 to AB 19 must be all 0s in order to access the RAM. There must be no unused address lines. A similar scheme could be done for the ROM, but all AB16-19 have to be 1.



5. **Instruction XCHG achieves a swap between the source operand and the destination operand. (Consult the 8086/88 user manual for more detailed information). Assume that the instruction XCHG does not exist in the 8086 instruction set. Write a sequence of instructions to duplicate the instruction XCHG AL,**

DL. Note that the values in all other registers (including AH and DH) should be their original values when your instruction sequence finishes.

PUSH CX Notes: 1. All stack operations are two bytes
MOV CL, AL 2. You can assume a memory location exists to act as the temporary store
MOV AL, DL
MOV DL, CL
POP CX

6. *Can the 8-bit input port at location 911 be accessed using direct port I/O addressing? Give an instruction sequence to copy the data from this port to register CL.*

Direct port I/O addressing requires that the address of the I/O port fits into a byte (from 0 to 255₁₀). Thus 911₁₀ is too large, and indirect port addressing (Using DX) must be used. Note that Direct/ Indirect I/O addressing is **INDEPENDENT** of the data size of the port (8-bit or 16-bit), and has nothing to do with memory addressing. The instruction should be:

MOV DX, 911 ; port address
IN AL, DX ; Input 8-bit from 911
MOV CL, AL ; copy data to CL

Also note that MOV CL, [911] is wrong! This refers to a memory location, not an I/O port.

7. *When a CALL is executed, how does the CPU know where to return? What is the difference between a FAR call and a NEAR call?*

The address of the instruction immediately following the CALL is stored on the stack. The last instruction of a called subroutine must be RET in order to the system to pop off the return address from the stack.

In the FAR CALL, both the CS and IP registers are saved on the stack, whereas in a NEAR CALL, only the IP register will be saved on the stack.

8. *Find the contents of the stack and stack pointer after the execution of the CALL instruction shown next: Assume that SS:1296 right before the execution of CALL and SUM is a NEAR procedure.*

CS:IP
2450:673A CALL SUM
2450:673D DEC AH

IP = 673D will be stored in the stack at 1295 and 1294, therefore SS:1295 = 67 and SS:1294 = 3D. And the stack pointer will point to 1294 then.

9. *Translate one of the following two quotes to 8-bit ASCII format (ignore the names and dates) – (Text book: Section 3.4):*

a. *“640K ought to be enough for anybody.” - Bill Gates, 1981*

36 34 30 4B 20 6F 75 67 68 74 20 74
6F 20 62 65 20 65 6E 6F 75 67 68 20
66 6F 72 20 61 6E 79 62 6F 64 79 2E

b. *“I think there is a world market for about 5 computers.”
– Thomas J. Watson, founder of IBM, 1943*

49 20 74 68 69 6E 6B 20 74 68 65 72 65 20 69 73
20 61 20 77 6F 72 6C 64 20 6D 61 72 6B 65 74 20
66 6F 72 20 61 62 6F 75 74 20 35 20 63 6F 6D 70
75 74 65 72 73 2E

10. Write your MUN student number and convert to its unpacked BCD binary equivalent (Text book: Section 3.4).

Simply match the corresponding number:

0 – 00H (00000000)	1 – 01H (00000001)	2 – 02H (00000010)
3 – 03H (00000011)	4 – 04H (00000100)	5 – 05H (00000101)
6 – 06H (00000110)	7 – 07H (00000111)	8 – 08H (00001000)
9 – 09H (00001001)		

eg. Student number : 9972563

decimal	9	9	7	2	5	6	3
BCD	1001	1001	0111	0010	0101	0110	0011
Or HEX	09H	09H	07H	02H	05H	06H	03H

11. Find the precise offset location in memory of each ASCII character or data in the following use a memory map (Text book: Section 3.4 for ASCII numbers):

ORG 20H

Data1 DB 73H, 2FH

DS:0020 73 DS:0021 2F

Data2 DB "737 3527"

DS:0022	37	DS:0023	33	DS:0024	37
DS:0025	20	DS:0026	33	DS:0027	35
DS:0028	32	DS:0029	37		

ORG 30H

Data3 DW 2560H, 101100010101B

DS:0030	60	DS:0031	25	DS:0032	15
DS:0033	0B				

ORG 40H

Data4 DD 25684FC4H

DS:0040	C4	DS:0041	4F	DS:0042	68
DS:0043	25				

Data5 DQ 7F5EC4527271FEH

DS:0044	FE	DS:0045	71	DS:0046	72
DS:0047	52	DS:0048	C4	DS:0049	5E
DS:004A	7F	DS:004B	00		

12. It is common practice to save all registers at the beginning of a subroutine. Assume that SP=1288H before a subroutine CALL. Show the contents of the stack pointer and the exact memory contents of the stack after PUSHF, for the following:

1132:0450 **CALL PROC1**

1132:0453 **INC BX**

.....

PROC1 PROC
PUSH AX
PUSH BX
PUSH CX
PUSH DX
PUSH SI
PUSH DI

...

...

PROC1 ENDP

When the procedure is called, IP, which points to the next instruction to be executed after the CALL, is saved on the stack since it is a NEAR procedure. After the CALL and all PUSH instructions have been executed, the stack is as follows with SP=127A

```

SS:127A    ←    DI
SS:127C    ←    SI
SS:127E    ←    DX
SS:1280    ←    CX
SS:1282    ←    BX
SS:1284    ←    AX          1285 = AH, 1284 = AL
SS:1286    ←    IP          1287 = 04, 1286 = 53
SS:1288

```

13. The following program adds four words and saves the result. The program contains some errors, fix the errors and make the program run correctly:

```

TITLE          PROBLEM PROGRAM
PAGE          60, 132
STSEG         SEGMENT
                DB 32 DUP(?)
STSEG         END           should be ENDS, Segment should end with ENDS
;-----
DTSEG         SEGMENT
DATA         DW 1234H, 3344H, 5FE2H, 85FAH
                ORG 10H
SUM          DW ?
DTSG         ENDS         Label should match
;-----
CDSEG         SEGMENT
START:       PROC FAR       Without : since it is a non-opcode generating instruction
                ASSUME CS:CDSEG, DS:DTSEG, SS:STSEG
                MOV DS, DTSEG   Could not move directly, should go through a register
                                Normally use MOV AX, DTSEG and MOV DS, AX
                MOV CS, 4       Could not move an immediate number to CS register
                MOV BX, 0
                MOV DI, OFFSET DATA
LOOP1        ADD BX, [DI]     Loop1 should have a colon :
                INC DI
                DEX BX         should be DEC
                JNZ LOOP1
                MOV SI, OFFSET RESULT
                MOV [SI], BX
                MOV AH, 4CH
                INT 21H
CDSEG:       ENDS         No : for CDSEG
START        ENDP         This line should be placed before the previous line
                END CDSEG     should be START

```

14. Write an Assembly Language Program that summarize eight unsigned byte numbers stored in memory and store the result back to the next memory location. The data segment can be defined as following:

```

DTSEG         SEGMENT
                Data         DB 23H, 34H, 32H, 45H, 1FH, 27H, 7FH, 90H
                Result       DW ?
DTSEG         ENDS

```

(Hint: 1. Use loop to write an efficient code. 2. Pay attention to how to handle the carry bit. You can use a 16-bit register to hold the result, and then you don't need to worry about the carry)

```

TITLE  ADDING 8 BYTES
PAGE   60, 132
STSEG  SEGMENT
       DB 64 DUP (?)
STSEG  ENDS
;-----
DTSEG  SEGMENT
       Data      DB    23H, 34H, 32H, 45H, 1FH, 27H, 7FH, 90H
       Result    DW    ?
DTSEG  ENDS
;-----
CDSEG  SEGMENT
MAIN   PROC  FAR
       ASSUME CS:CDSEG, DS:DTSEG, SS:STSEG

       MOV AX, DTSEG
       MOV DS, AX
       MOV CX, 8      ; COUNTER FOR 8 NUMBERS
       MOV SI, OFFSET DATA
       SUB AX, AX     ; CLEAR AX=0
REP:   SUB BX, BX     ; CLEAR BX=0
       MOV BL, [SI]
       ADD AX, BX
       INC SI
       LOOP REP
       MOV RESULT, AX ;SAVE RESULT BACK

       MOV  AH, 4CH
       INT 21H
MAIN   ENDP
CDSEG  ENDS
       END MAIN

```

Notes: This is just one example, should have many different ways to write this program.