

## Problem set 2

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Important: For all these problems, do not worry about efficiency. We will explore efficient approaches later. For now I want to focus on looking at problems as instances of more general sub problems that can be broken down.

Q0 We need to break a sequence of items of various weights into roughly equally weighted segments. The problem is represented as a sequence of  $n$  positive real weights  $w = [w_0, w_1, \dots, w_{n-1}]$ , a positive real number  $x$ , and a positive integer  $p$ . The goal is to break  $w$  into  $p$  segments such that no segment weighs more than  $x$ . For each segment, there is a penalty of the cube of the difference between the weight of the segment and  $x$ . If we represent the segments by an array of  $p + 1$  numbers  $k(0), k(1), \dots, k(p)$  where

$$0 = k(0) \leq k(1) \leq \dots \leq k(p-1) \leq k(p) = n,$$

then the segments are  $[w_{k(0)}, w_1, \dots, w_{k(1)-1}]$ ,  $[w_{k(1)}, w_{k(1)+1}, \dots, w_{k(2)-1}]$ ,  $\dots$  and  $[w_{k(p-1)}, w_{k(p-1)+1}, \dots, w_{k(p)-1}]$ . An acceptable solution has

$$w_{k(i)} + w_{k(i)+1} + \dots + w_{k(i+1)-1} \leq x$$

for each  $i \in \{0, \dots, p\}$ . An optimal solution minimizes the total penalty, which is

$$\sum_{i \in \{0, \dots, p\}} (x - w_{k(i)} - w_{k(i)+1} - \dots - w_{k(i+1)-1})^3.$$

Design a function to compute the cost (i.e., total penalty) of an optimal solution for inputs  $w$ ,  $x$ , and  $p$ . Your function should return  $\infty$  if there is no solution. Hint: For a given  $w$ ,  $x$ , and  $p$ , you can define subproblems defined by integers  $i$  and  $q$ , such that  $0 \leq i \leq n$  and  $0 \leq q \leq p$ ; subproblem  $(i, q)$  is to find the cost of an optimal way to split the first  $i$  items of  $w$  into  $q$  segments, each of which weighs less than  $x$ . The original problem is just the subproblem such that  $i = n$  and  $q = p$ .

- Specify using pre- and postconditions
- Write the function body

Q1 Suppose we represent a complex project by a simple directed acyclic graph  $G = (V, E)$ . Each vertex represents a milestone (including a start and finish milestone), while each edge represents a task to be done. Each task is associated with a time, which is the time it will take to complete. You need to find the total time of the longest path from the start to the finish.

- Specify subproblems using pre- and postconditions.
- Write the function body.

Q2 Given a set of items  $S$ , produce the set of all permutations of items in  $S$ . Try to solve this problem in two different ways.

Q3 Given two sequences, how many operations are needed to transform one into the other  
Each operation is one of

- Delete an item
- Add an item
- Replace one item with another

Example: This edit sequence has 7 operations. Is this minimal?

	midway upon the journey of our life
	<b>in</b> the midway of this our mortal life
insert "in"	
at 0	in midway upon the journey of our life
	in <b>the</b> midway of this our mortal life
insert "the"	
at 1	in the midway <b>upon</b> the journey of our life
	in the midway <b>of</b> this our mortal life
replace "upon" with "of"	
at 3	in the midway of <b>the</b> journey of our life
	in the midway of <b>this</b> our mortal life
replace "the" with "this"	
at 4	in the midway of this <b>journey</b> of our life
	in the midway of this our mortal life
delete "journey"	
at 5	in the midway of this <b>of</b> our life
	in the midway of this our mortal life
delete "of"	
at 5	in the midway of this our life
	in the midway of this our <b>mortal</b> life
insert "mortal"	
at 6	in the midway of this our mortal life
	in the midway of this our mortal life

- Identify the subproblems.
- Specify with pre- and postconditions.
- Design the body