

# Assignment 0 Discrete math.

## Advanced Computing Concepts for Engineering

Due January 11, 2018

This assignment will not be marked.

**Q0 Set notation.** Use the filter and/or map notations to concisely express the following sets

- (a) The set of all composite natural numbers.  $\{0, 4, 6, 8, 9, 10, 12, 14, 15, 16, 18, 20, \dots\}$ .
- (b) The set of all positive real numbers less than the square root of 2.

Don't use the square-root sign in your answer.

(c) The set of all straight lines in the Cartesian plane. (Consider a straight line to be a suitable subset of  $\mathbb{R} \times \mathbb{R}$ .)

### Q1 Paradoxical sets.

In class it was stated that a set is a collection of mathematical objects and that each set is itself a mathematical object. Let  $S$  be the set of all sets.

Presumably some sets contain themselves. E.g., since  $S$  is a set and  $S$  contains all sets, we can conclude that  $S \in S$ . On the other hand, it is clear that some sets do not contain themselves. For example  $\emptyset \notin \emptyset$ . So some sets contain themselves and some do not.

Let  $R$  be the set of all sets that do not contain themselves. I.e.  $R = \{x \in S \mid x \notin x\}$ .

(a) Prove that if  $R \notin R$  then  $R \in R$ . Can you conclude that it is not the case that  $R \notin R$ ?

(b) Prove that if  $R \in R$  then  $R \notin R$ . Can you conclude that it is not the case that  $R \in R$ ?

(c) What does this tell us about the set of all sets as a viable concept?

### Q2 Counting

Let  $S$  and  $T$  be finite sets. Let  $|S| = m$  and  $|T| = n$ .

(a) What is the size of  $|S \times T|$ .

(b) How many binary relations are there with  $S$  as source and  $T$  as target?

- (c) How many total functions are there with  $S$  as source and  $T$  as target?  
 (d) How many partial functions are there with  $S$  as source and  $T$  as target?

**Q3 Propositional logic.**

Using the laws that appear above them in the notes, prove the following distributivity laws

- (a)  $(p \wedge q \Rightarrow r) = ((p \Rightarrow r) \vee (q \Rightarrow r))$   
 (b)  $(p \Rightarrow q \wedge r) = ((p \Rightarrow q) \wedge (p \Rightarrow r))$

**Q4 Substitutions**

- (a) Underline all bound occurrences of variables in the following formulae. Circle all free occurrences of variables.

$$\{i \in \mathbb{N} \mid i < f(i) \cdot g(i)\}$$

$$(\forall x \in \mathbb{R} \cdot g(x) < f(y))$$

- (b) Make the following substitutions

$$\{i \in \mathbb{N} \mid i < f(i) \cdot g(i)\} [x, i, f : y, j, g]$$

$$(\forall x \in \mathbb{R} \cdot g(x) < f(y)) [y : y + 1]$$

$$(\forall x \in \mathbb{R} \cdot g(x) < f(y)) [y : x + 1]$$

**Q5 Quantifiers and sets**

Ranter is a social network in which users publish short messages called rants. Let  $U$  be the set of all users on the social network and let follows :  $U \times U \xrightarrow{\text{tot}} \mathbb{B}$  be a boolean function expressing that the first user follows the second.

- (a) Use quantifiers to say that following is irreflexive, i.e., that no one follows themselves.

(b) If, for some users  $a$  and  $b$ , follows( $a, b$ ), we say that  $a$  is a degree 1 follower of  $b$ . If, for some users  $a, b$ , and  $c$ ,  $a$  follows  $c$  and  $c$  follows  $b$ , we say that  $a$  is a degree 2 follower of  $b$ . And so on. Use quantifiers to say that  $a$  is a degree  $k$  follower of  $b$ . You may assume  $k > 0$ . [Hint: You may need to quantify over a function. Hint: check that the free variables of your expression are  $a, b$ , and  $k$ .]

- (c) Explain the meaning of the following expression in clear English

$$\forall x \in R \cdot \exists y \in S \cdot \text{follows}(x, y)$$

[Hint: check that the free variables of your English sentence are the same as the free variables of the expression.]

(d) Explain the meaning of the following expression in clear English

$$\exists x \in R \cdot \forall y \in S \cdot \text{follows}(x, y)$$

[Hint: check that the free variables of your English sentence are the same as the free variables of the expression.]

### Q6 Scheduling

Let  $S$  be a set of course sections. Let  $T$  be a set of times. Let  $R$  be a set of rooms.

Let  $at$  be a binary function  $at : S \times T \xrightarrow{\text{tot}} \mathbb{B}$ . The intended meaning of  $at(s, t)$  is that section  $s$  is scheduled at time  $t$ .

Let  $in$  be a binary function  $in : S \times T \times R \xrightarrow{\text{tot}} \mathbb{B}$ . The intended meaning of  $in(s, t, r)$  is that section  $s$  is scheduled to be in room  $r$  at time  $t$ .

Express the following statements using quantifiers

(a) All sections are scheduled at exactly 3 times, according to the  $at$  function.

(b) All sections are scheduled at times (according to  $at$ ) only when they are in a room.

(c) All sections are in a room, only at times they are scheduled to be at (according to  $at$ ).

(d) All sections are only scheduled to be in one room at a time.

(e) No room is in use by two (or more) sections at the same time.

Express the following sets using set notation and quantifiers.

(f) The set of all room/time pairs when 2 sections are both scheduled for the room at the time.

(g) The set of all section/time pairs which need a room.