

## Applying Predicate Logic to Software Documentation

Subroutines are often documented by specifying the following information

- Precondition: A boolean expression describing what should be true when the subroutine begins execution
- May Change: A list of variables whose values may be changed.
- Postcondition: A boolean expression describing what will be true when the subroutine completes execution
  - \* In the postcondition,  $v'$  represents the final value of variable  $v$ ,
  - \* while a plain  $v$  represents the initial value of the variable  $v$

**Example** a subroutine that searches an array  $A$  for a particular value  $x$  may be described by

```
void find( int A[N], int x, int &i )
// Precondition: exists j : {0,1,...,N-1}, A[j]==x
// May change: i
// Postcondition: A[i']==x
```

- The precondition says that the subroutine should only be called if there is an  $x$  somewhere in  $A$ .
- The postcondition says that the final value of  $i$  should index an element of  $A$  equal to  $x$ .

**Example** a subroutine that sorts an array of integers

```
void sort( int A[N] )
// Precondition: true
// May change: A
// Postcondition: (for all i : {1,2,...,N-1}, A'[i-1] <= A'[i])
//   and (for all x : Int, |\{i | A[i]==x\}| == |\{i | A'[i]==x\}| )
```

- The first line of the postcondition says that the array  $A$  is sorted at completion
- The second line says that its contents have been permuted, but not otherwise changed.

## Applying predicate logic to system specification

A “System” may be defined as

- an object that imposes a relationship on objects labeled as its inputs and outputs.

A “System model” is a boolean expression that describes the relationship the system imposes.

- The free variables of the model are the names of the inputs and outputs.
- Usually inputs and outputs are modelled as functions of time
- Often, but not always, a system model is a function (aka transform) from its inputs to its outputs.

Systems are often composed from subsystems

### An example:

(We take time to range over the natural numbers counting clock cycles.)

Thus  $x$ ,  $y$  and  $z$  range over functions from the natural numbers  $\mathbb{N}$  to the set  $\{T, F\}$

Consider a system consisting of a not-gate with input  $x$  and output  $y$

$$\text{Not}(x, y) \triangleq (\forall t, y(t) \leftrightarrow \neg x(t))$$

Consider a system consisting of a D-flip-flop with input  $x$  and output  $y$ .

$$\text{DFF}(x, y) \triangleq (\forall t, y(t+1) \leftrightarrow x(t))$$

We can compose these two systems in various ways.

Here are models of two

$$\text{NotThenDFF}(x, y) \triangleq \exists z, \text{Not}(x, z) \wedge \text{DFF}(z, y)$$

and

$$\text{DFFThenNot}(x, y) \triangleq \exists z, \text{DFF}(x, z) \wedge \text{Not}(z, y)$$

We can show (using predicate logic and other math) that these two composed systems have equivalent models.

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Although this example deals with digital systems, exactly the same ideas apply to analog systems.