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# Contracts for objects -- 0

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Clear Box Specification

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# Contracts for classes

- Now we extend the idea of contracts to classes.
- As an example, we consider a class for representing rational numbers.
- We use a simple data structure:

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# Rational

```
class Rational {  
    private double numerator ;  
    private double denominator ;  
  
    // requires d != 0.0  
    // ensures denominator' != 0.0  
    public Rational( double n, double d ) {  
        numerator = n ; denominator = d ; }  
}
```

# Rational

```
// requires denominator != 0.0  
// ensures result == numerator / denominator  
public double toDouble() {  
    return numerator / denominator ; }
```

- Does it make sense to require the client to ensure that the denominator is not 0 before calling toDouble?
- We should not force the client to reason in terms of the private fields of an object.
- To do so is contrary to the principles of information hiding and abstraction.

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- Objects are meant to represent *things*.
  - There are certain states of the objects that are sensible and certain states that --while representable by the fields-- should not be reachable. These states do not represent things.

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# Invariants

- It is the job of the implementer of a class (not its clients) to ensure that the objects of the class do not reach states that are not sensible.
- An object invariant is a description of the states that of an object that are sensible.
- We start again. This time we state the invariant

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# Invariants

```
class Rational {  
    // invariant denominator != 0.0  
    protected double numerator ;  
    protected double denominator ;  
  
    // requires d != 0.0  
    public Rational( double n, double d ) {  
        numerator = n ; denominator = d ; }  
    // ensures result == numerator / denominator  
    public double toDouble() {  
        return numerator / denominator ; }  
}
```

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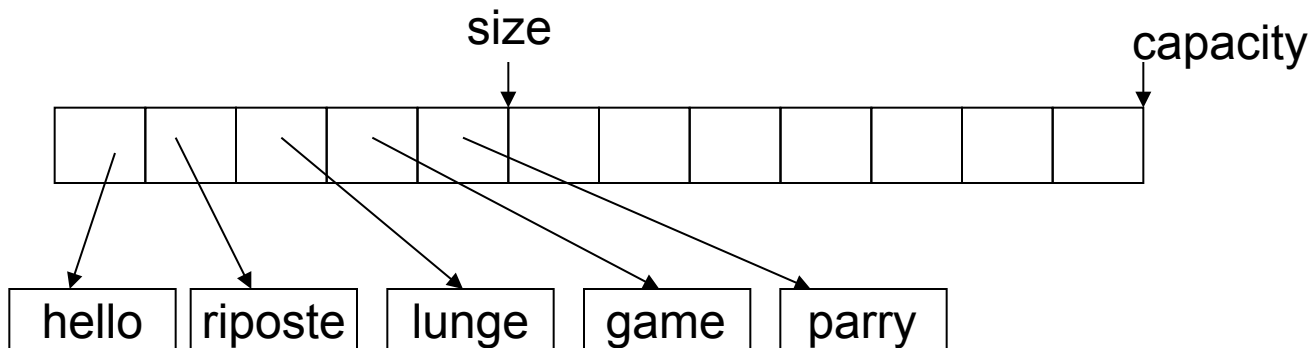
# Invariants

- The client coder does not need to think about the invariant.
- The implementer may assume that the invariant is true at the start of each method.
- But the implementer must also ensure that the each method and constructor of the class establishes the invariant at its end.
- Thus each method should preserve the invariant.



# Another example

- As a second example, we use a dictionary that creates and records an association between strings and small integers.
- We use a simple data structure:



# Data structure

```
class Dictionary {  
    public final static INIT_CAPACITY = 10 ;  
    protected int size = 0 ;  
    protected String[] a = new String[INIT_CAPACITY ] ;  
  
    // modifies size, a  
    // ensures size' == 0 and a' != null  
    public Dictionary() { ... }  
  
    // ensures result == size  
    public int getSize() { ... }  
  
    // requires a != null  
    // ensures result == a.length  
    public int getCapacity() { ... }
```

# getInt

```
// requires str != null  
// and a != null and 0 <= size and size <= a.length  
// and (for all i in {0,1,...,size-1}, a[i]!=null)  
// and (for all i,j in {0,1,...,size-1}, a[i]==a[j] implies i==j)  
// ensures  
// if( there is an i in {0,1,...,size-1}, str.equals(a[i]) )  
// then 0 <= result and result < size  
// and str.equals( a[result] ) )  
// else result == -1  
public int getInt(String str ) { ... }
```

# putString

```
// requires str != null and size < a.length
//   and a != null and 0 <= size and size <= a.length
//   and (for all i in {0,1,...,size-1}, a[i]!=null)
//   and (for all i,j in {0,1,...,size-1}, a[i]==a[j] implies i==j)
// modifies a[size], size
// ensures 0 <= result and result <= size'
//   and str.equals(a[result]) and (size' in {size, size+1})
//   and (for all i in {0,1,...,size-1}, a[i]'.equals(a[i])
//   and a != null and 0 <= size' and size' <= a'.length'
//   and (for all i in {0,1,...,size'-1}, a[i]'!=null)
//   and (for all i,j in {0,1,...,size'-1}, a[i]'==a[j]' implies i==j)
int putString( String str ) { ... }
```

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# Invariants

- Notice that certain facts about the fields are required by almost all methods.
- Thus these facts must be established by each constructor and *preserved* by each method
- These facts essentially define what it means for the state of the object to be *sensible*.

# Invariants

- In this example, we require
  - That `a` points to an array:
    - `a != null`
  - That `size` is a valid index or equals the capacity:
    - `0 <= size` and `size <= a.length`
  - That the first `size` items of the array are not null:
    - **(for all  $i$  in  $\{0, 1, \dots, \text{size}-1\}$ ,  $a[i] \neq \text{null}$ )**
  - That the first `size` items of `a` be unique:
    - **(for all  $i, j$  in  $\{0, 1, \dots, \text{size}-1\}$ ,  $a[i] == a[j]$  implies  $i == j$ )**
  - If any of these “facts” is false, then the data structure is corrupt.

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# Invariants

- We call these facts the *object invariant* (sometimes called *class invariant*)
- The object invariant must be ensured by each constructor and each method of the class.
- The invariant may thus be assumed at the start of each method.

# Rewriting the class

- Now we rewrite the Dictionary class, factoring out the invariant.

```
class Dictionary {  
    public final static INIT_CAPACITY = 10 ;  
    protected int size = 0 ;  
    protected String[] a = new String[ INIT_CAPACITY ] ;  
    // invariant a != null  
    // invariant 0 <= size and size <= a.length  
    // invariant (for all i in {0,1,...,size-1}, a[i]!=null)  
    // invariant (for all i, j in {0,1,...,size-1}, a[i]==a[j] implies i==j)
```



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# Rewriting the class

```
// modifies size, a  
// ensures size == 0  
public Dictionary() { ... }  
  
// ensures result == size  
public int getSize() { ... }  
  
// ensures result == a.length  
public int getCapacity() { ... }
```

# Rewriting the class

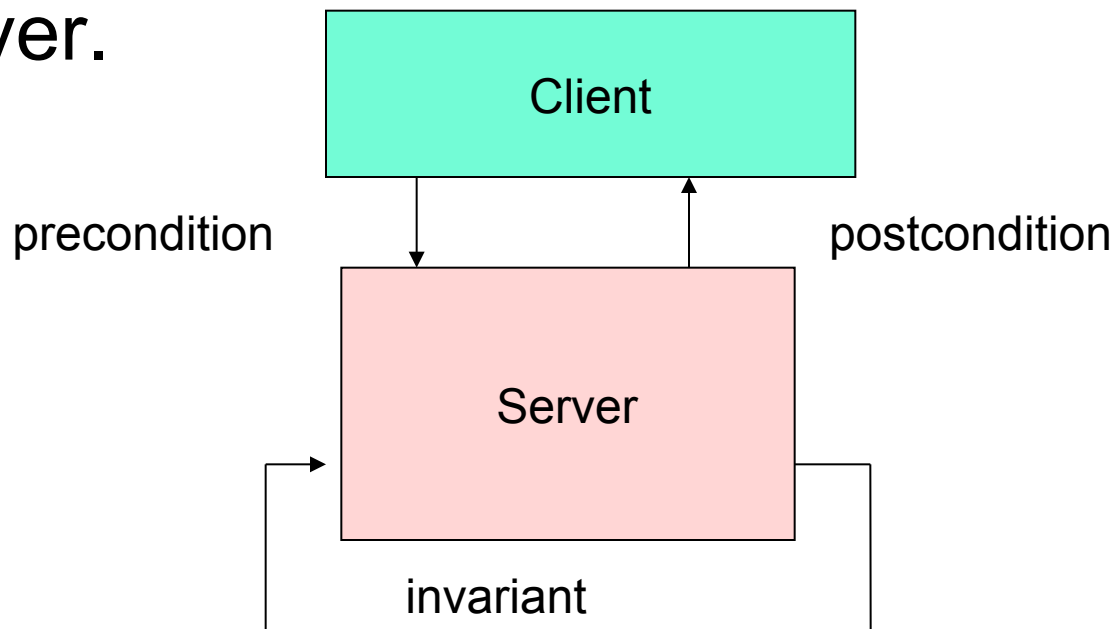
```
// requires str != null  
// ensures  
// if( there is an i in {0,1,...,size-1}, str.equals(a[i]) )  
// then 0 <= result and result < size  
//      and str.equals( a[result] ) )  
// else result == -1  
public int getInt(String str ) { ... }
```

# Rewriting the class

```
// requires str != null and size < a.length  
// modifies a[size], size  
// ensures 0 <= result and result <= size'  
// and str.equals(a[result])' and (size' in {size, size+1})  
// and (for all i in {0,1,...,size-1}, a[i].equals(a[i])  
int putString( String str ) { ... }
```

# Summary

- Note that the precondition now contains only things that the client actually has control over.



Flow of obligations

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# Invariants and defensive checks

- We can typically write the invariant as a method that is called at the end of each constructor and mutator (method that changes state). The check can be partial or full.
- To be extra careful, also call it at the start of each method.

```
protected void invariant () {  
    assert a != null ;  
    assert 0 <= size && size <= a.length ; ... }  

```

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# Invariants and callbacks

- As mentioned, it is ok for the invariant to become untrue during the execution of a method, as long as it is restored by the end.
- Of course the invariant must be true also before any call that might cause a method invocation on the same object.
- In particular you have to be careful about calling other objects that might call back

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E.g.

```
void someMutator() {  
    ...make some changes...  
    invariant() ; // invariant should be true here  
    notifyAllObservers() ;  
    ... do something else...  
    invariant() ;  
}
```

# Invariants and shared objects

- Recall that in concurrent programming we should ensure that *shared objects* are never “owned” (aka “occupied”) by more than one thread at a time.
- The invariant of a shared object should be true whenever no thread owns it.
- It may be assumed at the start of synchronized methods.
- It should be true on return from synchronized methods.
- It should be true before any call to `wait()`.
- It may be assumed after any call to `wait()`.
- I.e. it is both a pre- and a postcondition of `wait()`.