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# Design by Contract

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# Contracts

- Alice (client) hires Bob (server) to fix her car.
- They make a contract.
  - Alice agrees to give Bob \$100 in advance
  - Bob agrees that when he is done, the car will be in good working order

# Contracts

	Obligation	Benefit
Alice (client)	Must pay \$100	Has working car
Bob (server)	Must fix car	Has \$100 to buy materials

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# Contracts in programming

- Consider a method **tan** to compute the tangent of an angle between 0 and 89 degrees.
- Alice will write the client.
- Bob will implement **tan**.
- Contract
  - Syntactic signature: **double** tan( **double** x)
  - Alice agrees to send, as argument, a value between 0 and 89.
  - Bob agrees that the result will be equal to the tangent to at least three decimal places.

# Contracts in programming

	Obligation	Benefit
Alice (client)	Must supply an argument in range [0,89]	Result equals the tan of the argument to 3 decimal places
Bob (implementer)	Must ensure the result equals the tan of the argument to 3 decimal places	Argument in range [0,89]

# Contracts in programming

From the point of view of the implementer:

- the clients obligation is what is **required**.
- the implementer's obligation is what the implementation **ensures**.
- We document procedures as follows

```
/**
```

```
* requires boolean expression
```

```
* ensures boolean expression
```

```
*/
```

“requires clause” -- also called  
“precondition”

“ensures clause” – also called  
“postcondition”

# Contracts in programming

We use “result” to represent the result of an invocation.

## ■ For example:

```
class DegMath {  
    /** requires 0 <= x && x <= 89  
    * ensures result == tan(x) to three decimal places,  
    * where tan is the mathematical tangent function in  
    * terms of degrees. */  
    static double tan( double x )  
}
```

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## If the client breaks the contract

- Given this contract, the client can not make any assumptions about what a call such as `DegMath.tan( 90.0 )` or `DegMath.tan( -1.0 )` might do.
- *The client is obligated to ensure that the expression in the requires clause is true at the **start** of the invocation.*

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# If the client respects the contract

- For the implementation to be correct, the implementer must ensure that  
    `DegMath.tan( 25 )`  
equals the mathematically correct value to 3 decimal places.
- *The implementer is obligated to ensure that the “postcondition” is true, but only in those cases where the “precondition” is true.*



# Final values

## ■ Example

```
class Point {  
    double x, y ;  
  
    /** requires true  
    * modifies x, y  
    * ensures x'==0.0 && y'==0.0  
    Point() { ... }  
  
    /** requires true  
    * modifies x, y  
    * ensures x' == x + deltaX && y' == y + deltaY  
    */  
    void move( double deltaX, double deltaY ) {...}  
    ...  
}
```

# Omitting the requires clause

- In this example, there is no obligation on the caller beyond the syntactic signature. We can omit the requires clauses

```
class Point {  
    double x, y ;  
  
    /** modifies x, y  
    * ensures x'==0.0 && y'==0.0  
    Point() { ... }  
  
    /** modifies x, y  
    * ensures x' == x + deltaX && y' == y + deltaY  
    */  
    void move( double deltaX, double deltaY ) {...}  
    ...  
}
```

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# Framing

- The **modifies** clause is used to indicate which variables may be changed by a method.

```
class Point {  
    double x, y ;  
  
    /** modifies x  
     * ensures  $x' == x + \text{deltaX}$   
     */  
    void moveLeft( double deltaX ) {...}  
    ...  
}
```

# Example

## ■ Partitioning a list

```
/** requires a != null and a.length > 0
 * modifies a[*]
 * ensures a[*]' is a permutation of a[*] and
 * 0 <= result and result < a.length and
 * (for all i in {0,1,...,a.length-1}
 *     (i < result implies a[i]' <= a[result]') and
 *     (i > result implies a[i]' >= a[result]') )
 **/
int partition( double[] a )
```

## Further reading

- The paper that introduced the term “design by contract” was

Meyer, Bertrand. "Applying 'design by contract'." *Computer* 25, no. 10 (1992): 40-51.

- The ideas date back to the late 60s and 70s. For example, the Euclid programming language, designed in 1977, had support for pre- and postconditions.
- Meyer’s paper was important for applying the ideas to object-oriented programming.